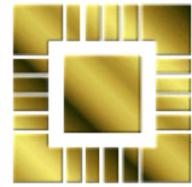


**CM214-COMP2008  
Data Communications and Networks**

**Network Security - 1**

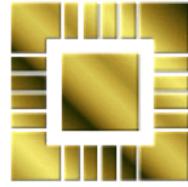
Karl R. Wilcox  
[krw@ecs.soton.ac.uk](mailto:krw@ecs.soton.ac.uk)



# Objectives



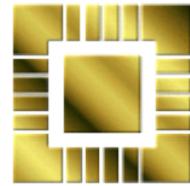
- To look at definitions of a “secure” network
  - To consider various aspects of security
  - To look at tools and techniques
- (Peterson & Davie, Sections 8.1-3)



# Background Reading



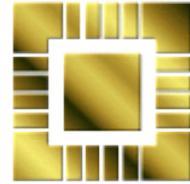
- “Security Complete”
  - Sybex, ≈£13, paperback, ≈900 pages
  - Comprehensive, practical, dry
- “Cryptonomicon”
  - Neal Stephenson, ≈ £10 paperback, ≈ 900 pages
  - Practical, interesting + good study



# What is “Security”? - 1



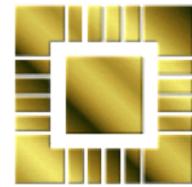
- That content of data in a network transfer remains private between sender & receiver(s)
  - Encryption
- That content of data in a network transfer has not been altered during transmission
  - Message Integrity



# What is "Security"? – 2



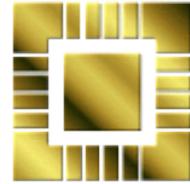
- That one or all parties in a network transfer can be assured of the identities of the other parties
  - Authentication
- That it can be proven that particular data was transferred between particular parties at a particular time
  - Non-repudiation



# An Alternative “Secure” Network



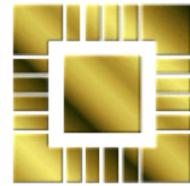
- All network transfers between any parties are untraceable with complete deniability(!)
  - E.g. A peer-to-peer file swapping network



# Security is NOT



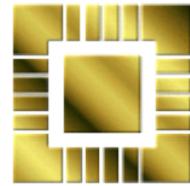
- Encryption
  - Although encryption plays a part
- Security is not the same as resilience
  - Although often confused / misused
- Resilience is resistance to failure
  - (including that caused by a deliberate attack)



# Security Is...



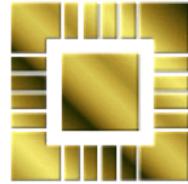
- Different things to different people
- A system built of components
  - We will look at tools / technologies
  - None are the single answer to “security”
- System is as secure as the weakest link
  - e-mail encrypted with 4096 bit RSA is not secure if it can be read over your shoulder



# Building a Secure System



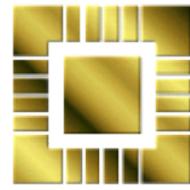
- Determine the required features
- Choose the tools
  - E.g. DES, IDEA, RSA, MD5
- Choose the techniques
  - E.g. Authentication, digital signatures, key distribution
- Or the applications
  - SSL, SSH, PGP



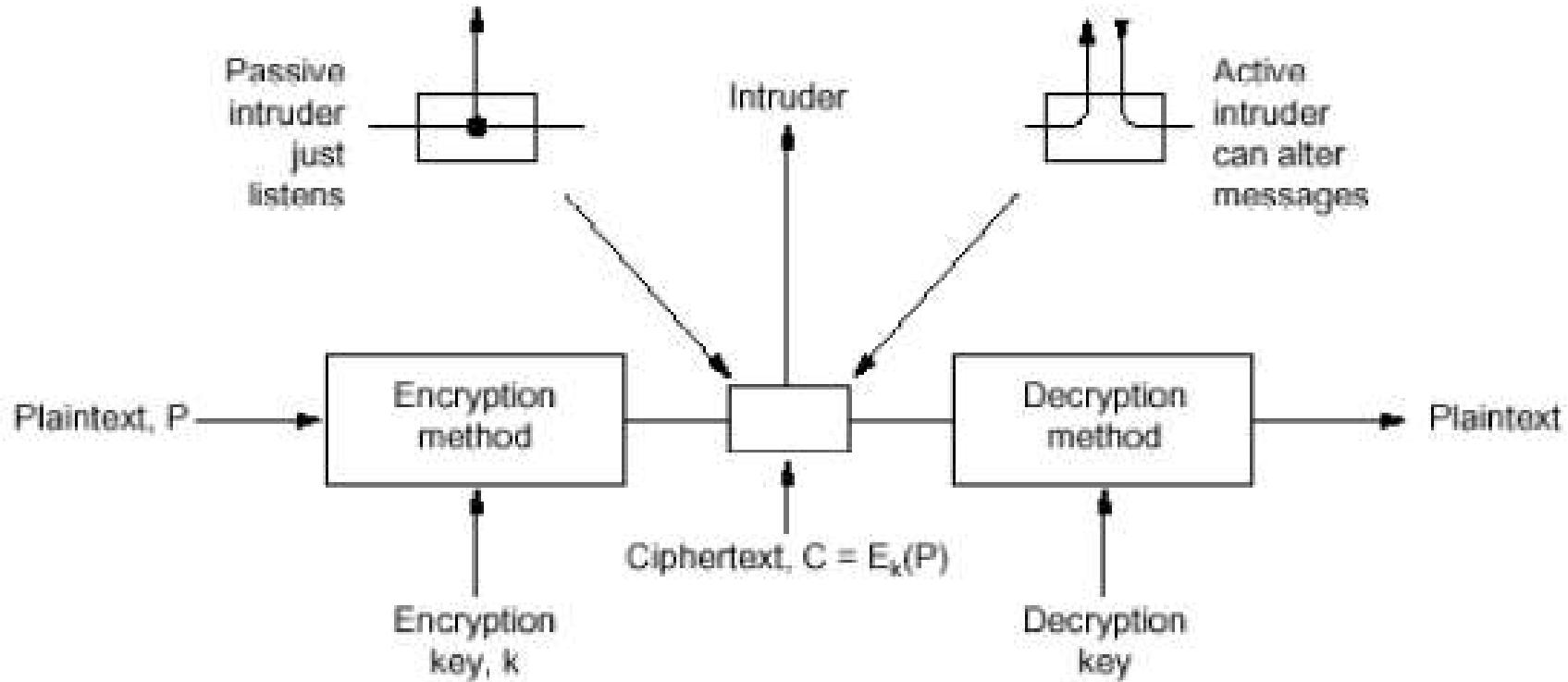
# Security Threats



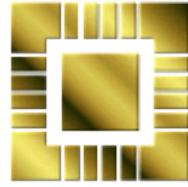
- Eavesdropping
- Impersonation
  - Person
  - Address
  - Computer
- Message duplication
- Key interception
- Cracking
  - Known plaintext
  - Brute force
- Line security
- Social engineering



# Security Model



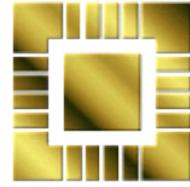
By convention, the parties in a transfer are known as “Alice” and “Bob”



# Traditional Cryptography



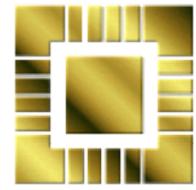
- Substitution / transposition cyphers
- Little use for network security
- Too open to attack
  - Known plaintext
  - Letter frequency
  - Brute force



# One Time Pads



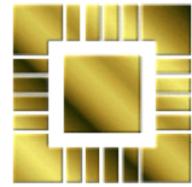
- Traditional, totally secure method
- Plain text of  $P$  bits
- Key  $K_B$  of at least  $B$  bits
- Coded message is  $C = K_B \ xor \ P$
- Decode by applying same key



# One Time Pad Requirements



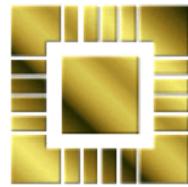
- Totally secure & unbreakable if:
  - Keys are genuinely random
  - The pad remains secret
  - It is never re-used
- Sounds ideal (simple, unbreakable)
  - Key distribution is the problem
  - Need to securely share as much key information as message information



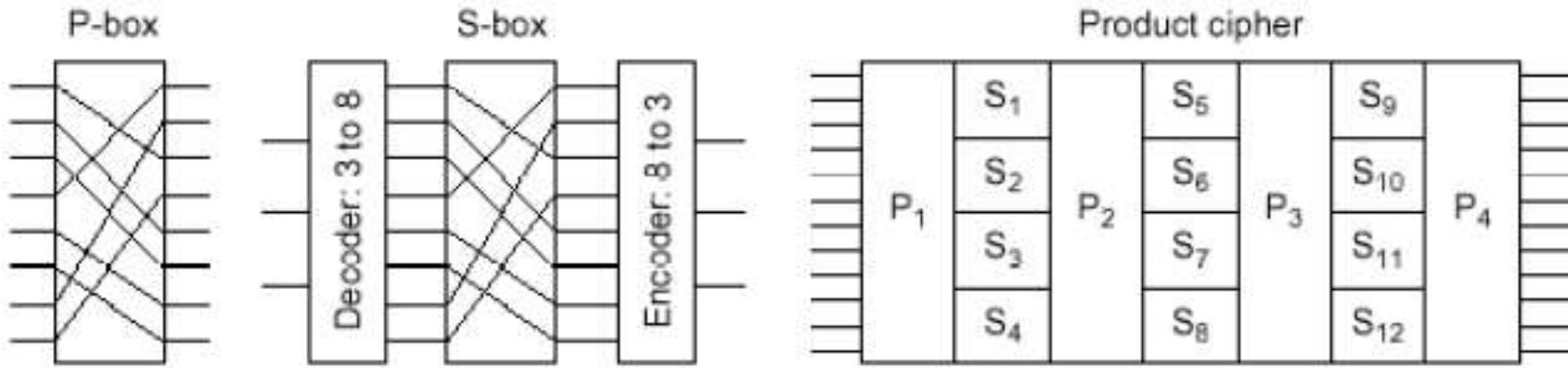
# Network Cryptography



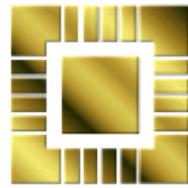
- Uses simpler keys (shorter than plaintext)
- Complex algorithms
  - Lots of iterations
  - Permuting bits
  - Substituting bit sequences



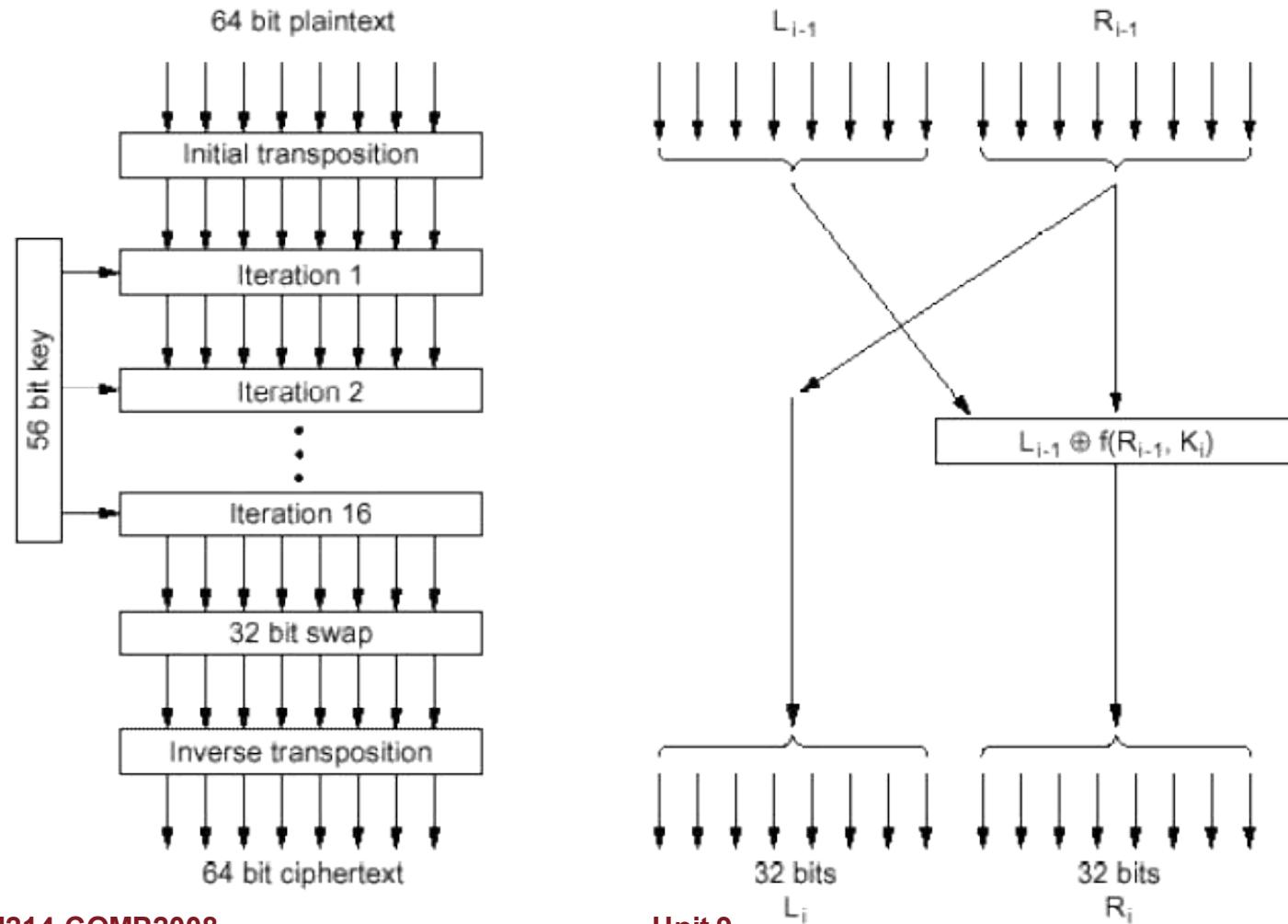
# Cryptography Elements

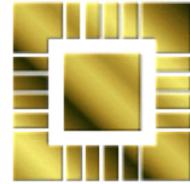


- Easy to implement in hardware
- Reasonably easy to implement in software
  - But computationally intensive



# Data Encryption Standard

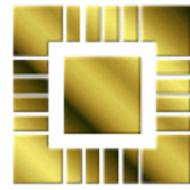




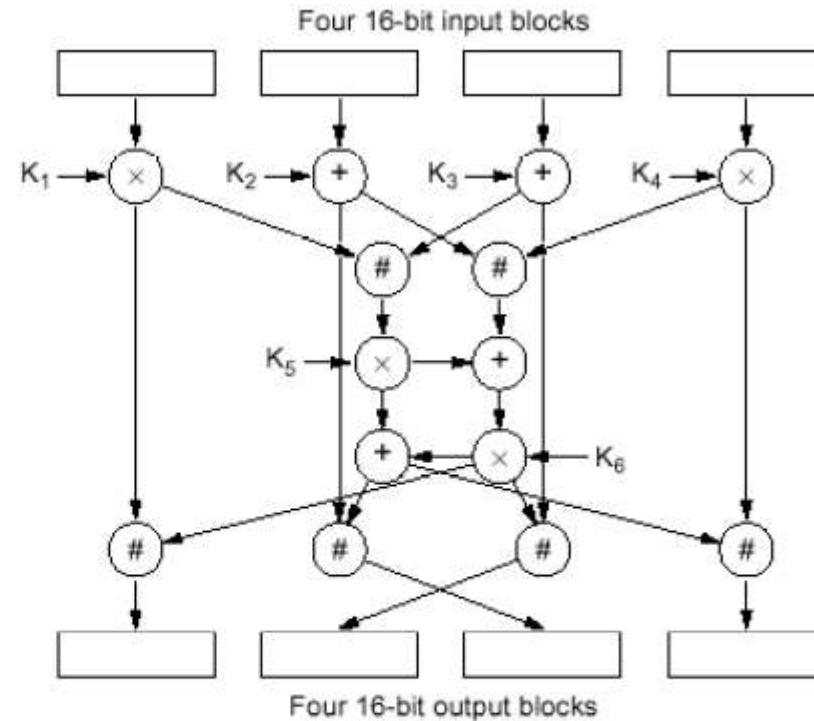
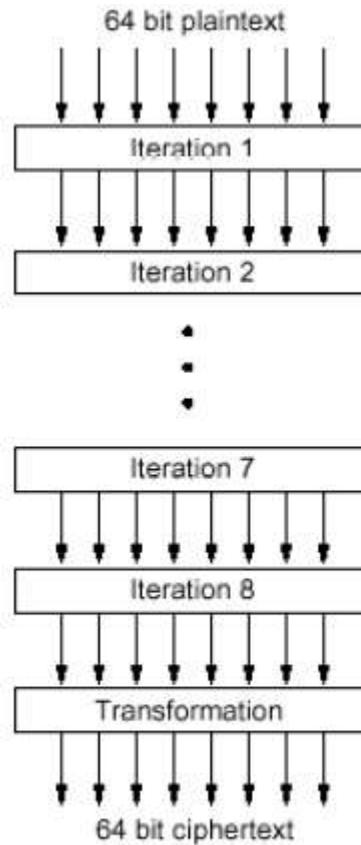
# How Secure is DES?



- Can be broken
  - By brute force in a “reasonable” time
    - 1999 \$250K computer 22 hours
  - Even faster with some known plain text
- Triple DES
  - Apply 3 times, 2 keys, 112 encryption
    - Cannot be broken in “reasonable” time



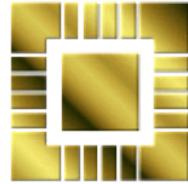
# International Data Encryption Standard



16-Bit addition modulo  $2^{16}$

16-Bit multiplication modulo  $2^{16} + 1$

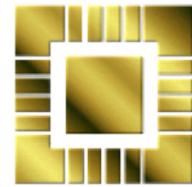
16-Bit EXCLUSIVE OR



# Encryption Keys



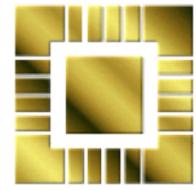
- Private keys (symmetric encryption)
  - Known only to involved parties
  - How to distribute?
- Public keys (asymmetric encryption)
  - One key is public (published)
  - One key remains private
  - Needs 2 pairs of keys for two-way communication



# Public Key Algorithms



- Use two complementary keys  $K_E$  &  $K_D$
- $D(E(P)) = P$
- Need to ensure:
  - $D$  not easily generated from  $E$
  - $E$  cannot be broken by plain text attack
- $E$  is public,  $D$  must remain private
- Any algorithm meeting these criteria will work



# Rivest, Shamir, Alderman (RSA) Algorithm



1. Take two (100+ digit) prime numbers  $p$  and  $q$
2. Calculate  $n = pq$ ,  $z = (p - 1)(q - 1)$
3. Choose  $d$  relatively prime to  $z$
4. Find  $e$  such that  $ed = 1 \bmod z$

To encode:

$$C = P^{**} e \bmod n$$

To decode:

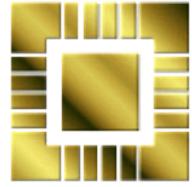
$$P = C^{**} d \bmod n$$



# RSA Features



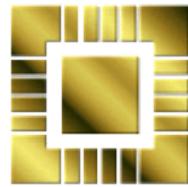
- Publish  $e$  and  $n$
- Keep  $d$  private
- Note, if we can factorise  $n$  we can break encryption
  - Fortunately  $n$  is very, very large



# RSA Example



- Pick two prime numbers,  $p = 3$ ,  $q = 11$
- This gives  $n = 33$ ,  $z = 20$
- Choose a number  $d$  with no common divisors with 20, say 7
- We want to find another number  $e$  that when multiplied by  $d$  divided modulo 20 leaves remainder 1
  - $e = 3$

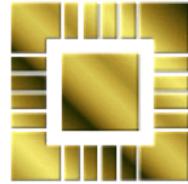


# RSA Calculations



Plaintext (P)		Ciphertext (C)		After decryption	
Symbolic	Numeric	$\underline{P^3}$	$\underline{P^3 \pmod{33}}$	$\underline{C^7}$	$\underline{C^7 \pmod{33}}$
S	19	6859	28	13492928512	19
U	21	9261	21	1801088541	21
Z	26	17576	20	1280000000	26
A	01	1	1	1	1
N	14	2744	5	78125	14
N	14	2744	5	78125	14
E	05	125	26	8031810176	5

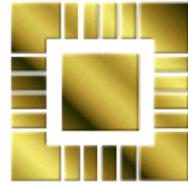
- Even with “toy” example consider large numbers involved
  - Infinite precision integer arithmetic



# Message Digests



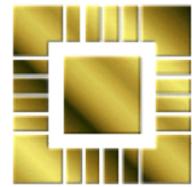
- Encryption is computationally intensive
- May not be primary aim
  - May wish to ensure message not tampered with in transit
- Need a “one-way” function between plain text and a (shorter) bit string
  - i.e. can generate bit string from plain text but not vice versa
  - Still need to encrypt digest



# Signatures



- An encrypted message digest also acts as a digital signature
  - The message is sent in plain text
  - Only the digest is encrypted
  - Only this message could generate the digest
  - Only the sender could have encrypted the digest



# Summary



- Encryption has more than one purpose
- And more than one way of implementation
- Self test – You wish to securely transmit a very long message
  - Should you compress or encrypt first?