



CM214-COMP2008
Data Communications and Networks
Network Security - 1

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Objectives



- To look at definitions of a “secure” network
 - To consider various aspects of security
 - To look at tools and techniques
- (Peterson & Davie, Sections 8.1-3)



Background Reading



- “Security Complete”
 - Sybex, ≈£13, paperback, ≈900 pages
 - Comprehensive, practical, dry
- “Cryptonomicon”
 - Neal Stephenson, ≈ £10 paperback, ≈ 900 pages
 - Practical, interesting + good stody



What is “Security”? - 1



- That content of data in a network transfer remains private between sender & receiver(s)
 - Encryption
- That content of data in a network transfer has been altered during transmission
 - Message Integrity



What is "Security"? – 2



- That one or all parties in a network transfer can be assured of the identities of the other parties
 - Authentication
- That it can be proven that particular data was transferred between particular parties at a particular time
 - Non-repudiation



An Alternative “Secure” Network



- All network transfers between any parties are untraceable with complete deniability(!)
 - E.g. A peer-to-peer file swapping network



Security is NOT



- Encryption
 - Although encryption plays a part
- Security is not the same as resilience
 - Although often confused / misused
- Resilience is resistance to failure
 - (including that caused by a deliberate attack)



Security Is...



- Different things to different people
- A system built of components
 - We will look at tools / technologies
 - None are the single answer to “security”
- System is as secure the weakest link
 - e-mail encrypted with 4096 bit RSA is not secure if it can be read over your shoulder



Building a Secure System



- Determine the required features
- Choose the tools
 - E.g. DES, IDEA, RSA, MD5
- Choose the techniques
 - E.g. Authentication, digital signatures, key distribution
- Or the applications
 - SSL, SSH, PGP



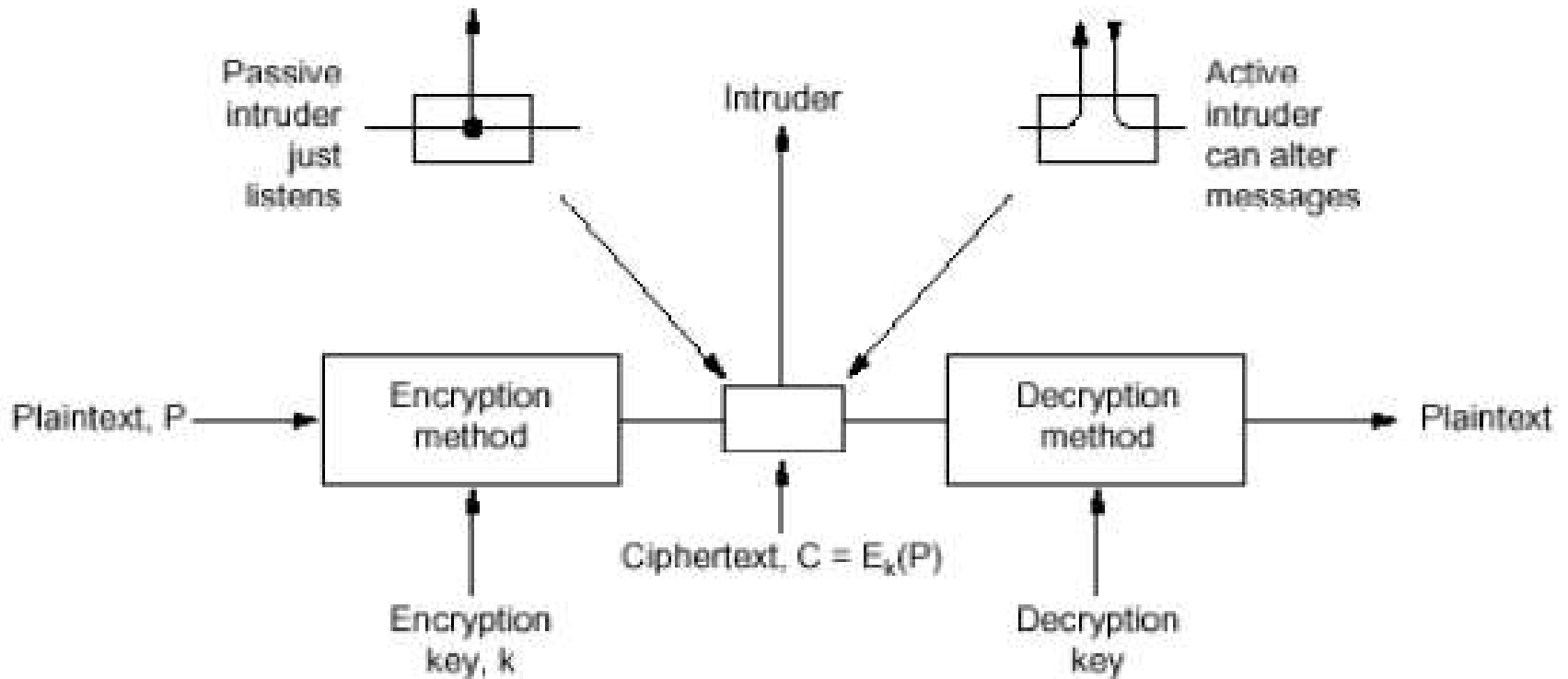
Security Threats



- Eavesdropping
- Impersonation
 - Person
 - Address
 - Computer
- Message duplication
- Key interception
- Cracking
 - Known plaintext
 - Brute force
- Line security
- Social engineering



Security Model



By convention, the parties in a transfer are known as “Alice” and “Bob”



Traditional Cryptography



- Substitution / transposition cyphers
- Little use for network security
- Too open to attack
 - Known plaintext
 - Letter frequency
 - Brute force



One Time Pads



- Traditional, totally secure method
- Plain text of P bits
- Key K_B of at least B bits
- Coded message is $C = K_B \text{ xor } P$
- Decode by applying same key



One Time Pad Requirements



- Totally secure & unbreakable if:
 - Keys are genuinely random
 - The pad remains secret
 - It is never re-used
- Sounds ideal (simple, unbreakable)
 - Key distribution is the problem
 - Need to securely share as much key information as message information



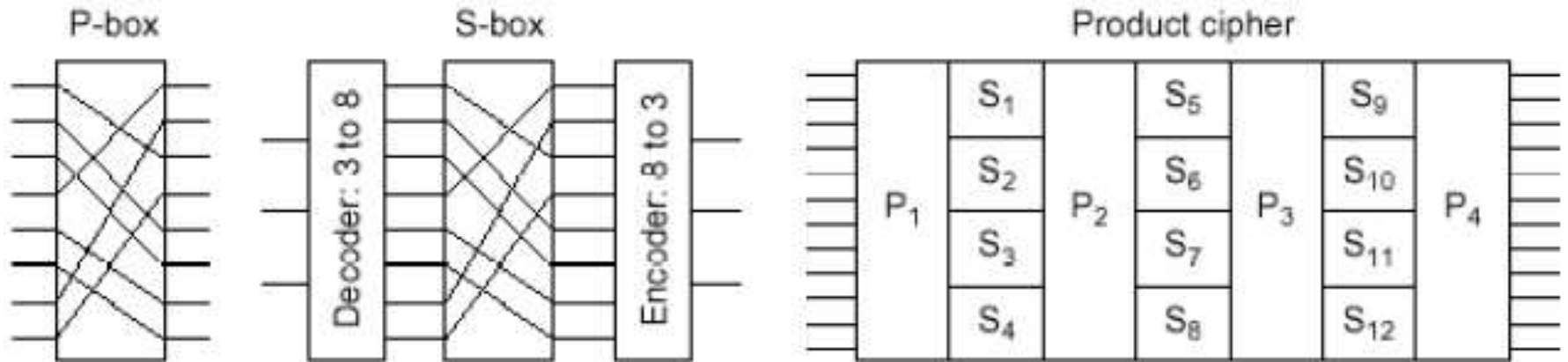
Network Cryptography



- Uses simpler keys (shorter than plaintext)
- Complex algorithms
 - Lots of iterations
 - Permuting bits
 - Substituting bit sequences



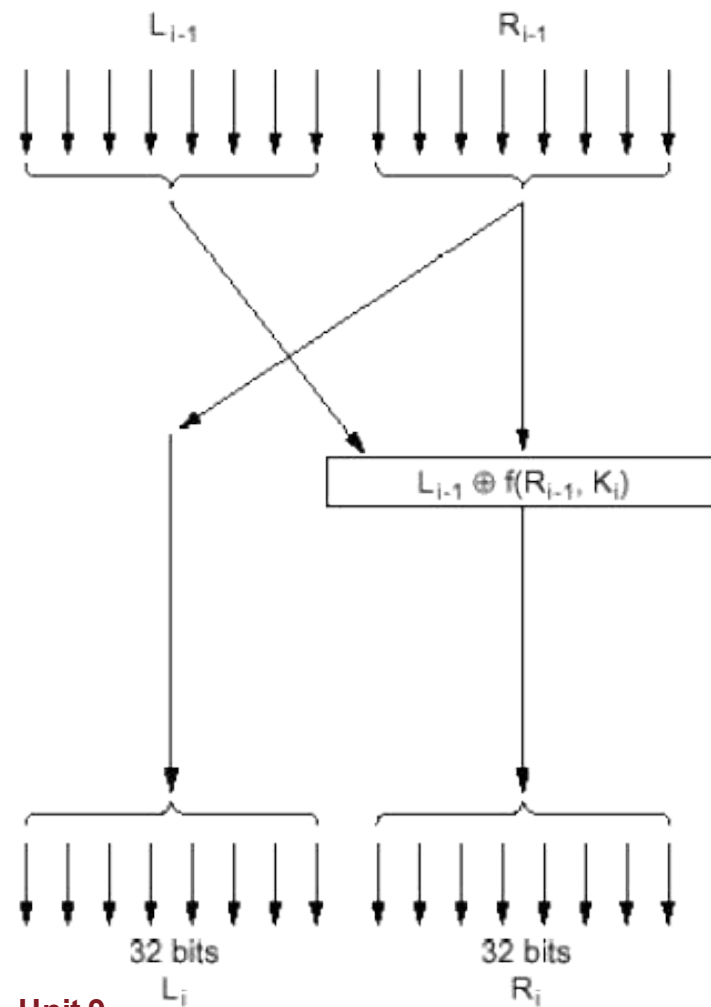
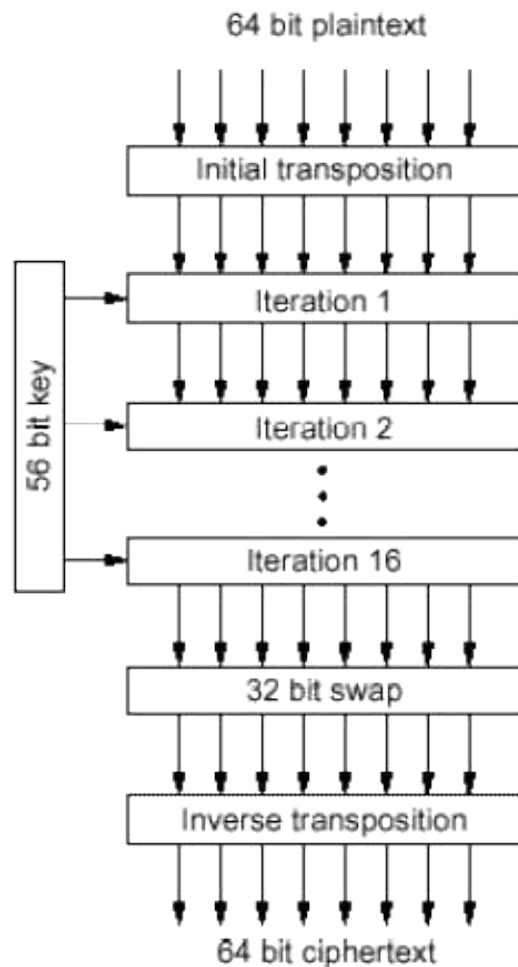
Cryptography Elements



- Easy to implement in hardware
- Reasonably easy to implement in software
 - But computationally intensive



Data Encryption Standard





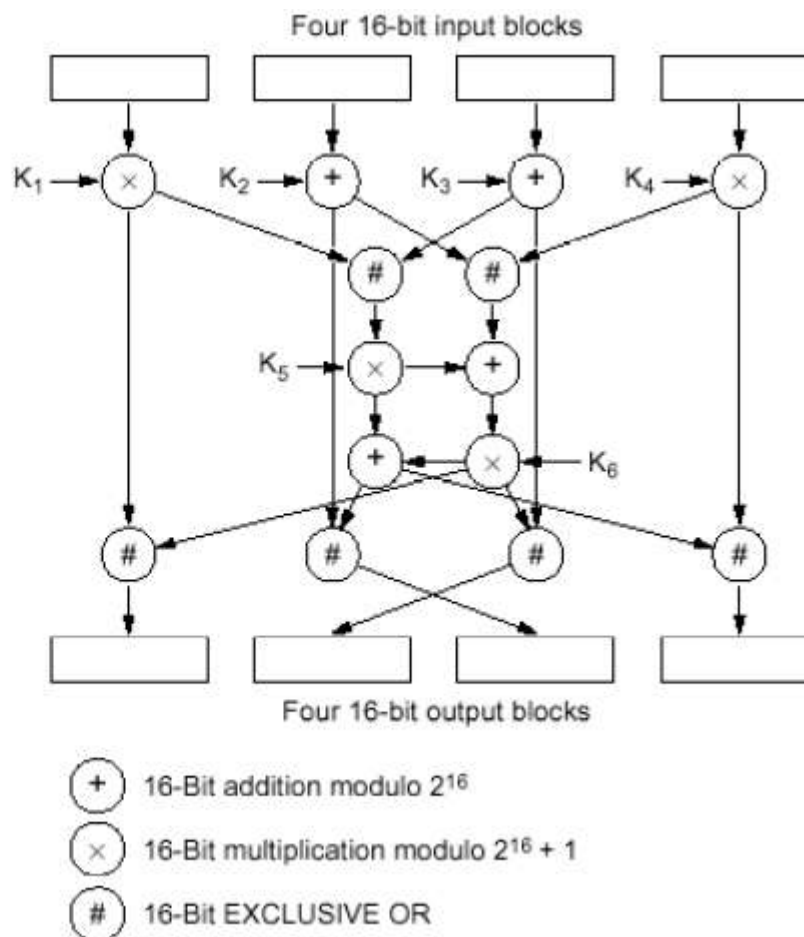
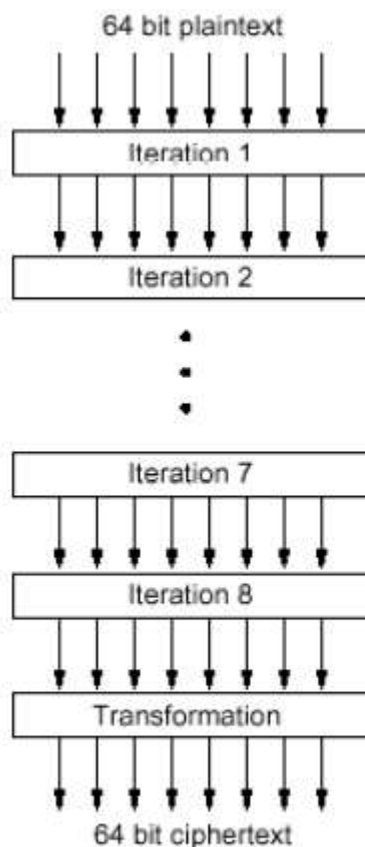
How Secure is DES?



- Can be broken
 - By brute force in a “reasonable” time
 - 1999 \$250K computer 22 hours
 - Even faster with some known plain text
- Triple DES
 - Apply 3 times, 2 keys, 112 encryption
 - Cannot be broken in “reasonable” time



International Data Encryption Standard





Encryption Keys



- Private keys (symmetric encryption)
 - Known only to involved parties
 - How to distribute?
- Public keys (asymmetric encryption)
 - One key is public (published)
 - One key remains private
 - Needs 2 pairs of keys for two-way communication



Public Key Algorithms



- Use two complementary keys K_E & K_D
- $D(E(P)) = P$
- Need to ensure:
 - D not easily generated from E
 - E cannot be broken by plain text attack
- E is public, D **must** remain private
- Any algorithm meeting these criteria will work



Rivest, Shamir, Alderman (RSA) Algorithm



1. Take two (100+ digit) prime numbers p and q
2. Calculate $n = pq$, $z = (p - 1)(q - 1)$
3. Choose d relatively prime to z
4. Find e such that $ed = 1 \bmod z$

To encode:

$$C = P ** e \bmod n$$

To decode:

$$P = C ** d \bmod n$$



RSA Features



- Publish e and n
- Keep d private
- Note, if we can factorise n we can break encryption
 - Fortunately n is very, very large



RSA Example



- Pick two prime numbers, $p = 3$, $q = 11$
- This gives $n = 33$, $z = 20$
- Choose a number d with no common divisors with 20, say 7
- We want to find another number e that when multiplied by d divided modulo 20 leaves remainder 1

$$- e = 3$$



RSA Calculations



Plaintext (P)		Ciphertext (C)			After decryption	
Symbolic	Numeric	P^3	$P^3 \pmod{33}$	C^7	$C^7 \pmod{33}$	Symbolic
S	19	6859	28	13492928512	19	S
U	21	9261	21	1801088541	21	U
Z	26	17576	20	1280000000	26	Z
A	01	1	1	1	1	A
N	14	2744	5	78125	14	N
N	14	2744	5	78125	14	N
E	05	125	26	8031810176	5	E

Sender's computation
Receiver's computation

- Even with “toy” example consider large numbers involved
 - Infinite precision integer arithmetic



Message Digests



- Encryption is computationally intensive
- May not be primary aim
 - May wish to ensure message not tampered with in transit
- Need a “one-way” function between plain text and a (shorter) bit string
 - i.e. can generate bit string from plain text but not vice versa
 - Still need to encrypt digest



Signatures



- An encrypted message digest also acts as a digital signature
 - The message is sent in plain text
 - Only the digest is encrypted
 - Only this message could generate the digest
 - Only the sender could have encrypted the digest



Summary



- Encryption has more than one purpose
- And more than one way of implementation
- Self test – You wish to securely transmit a very long message
 - Should you compress or encrypt first?